# Foundations of Data-Flow Analysis and <br> Constant Propagation 

## VYPe

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December 4, 2012

## Content

- Data-Flow Analysis
- Constant Propagation


## What is Data-Flow Analysis?

- Data flow analysis (DFA) is a special form of static analysis
- General steps:

1. Transform program into a Control flow graph (CFG)
2. Choose a property to inspect, eg. live variables, available expressions, constants. (property defines flow functions as in "flow of information")
3. Repeatedly apply flow functions to the CFG until a solution is found (maximum fix-point)

## Basic Blocks \& Control flow graph

- Basic block ( BB ) is a sequence of statements that is executed as a whole (atomically)
- Always entered via first statement (noone jumps inside the BB)
- The whole BB is executed (contains no jump instruction)
- Atomic execution of the BB ends by last statement (may be a jump, label etc.)
- Control flow graph (CFG) is a directed graph
- Equivalent to the original program
- nodes $\approx \mathrm{BBs}$
- edges $\approx$ transfers of control (eg. jumps) between BBs


## Sample program (spoiler: available expressions)

```
int a[10];
int b[10];
for(int i = 0 ; i < 10 ; i++) {
    if (i % 2 == 0)
        sum1 += a[i] * b[i]; // indexed access
    else
        sum2 += a[i] + b[i]; // indexed access
    if (a[i] > b[i]) // indexed access, again
        i++;
}
```


## Available expressions: redundancies elimination



## How to find available expressions (or other things)?

- A generic mathematical framework exists (magic with lattices)
- How it works
- Pick a property (eg. available expressions, live variables)
- Define flow functions (how to combine information from adjacent BBs)
- Attach an input and output set $I n_{b}$, Out $t_{b}$ to every basic block (set of available expressions, set of live variables etc.)
- Repeatedly recompute $I n_{b}, O u t_{b}$ sets of all BBs

- Fix-point is found, we are done


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## Constant propagation

If on every path leading to the point $p$ the expression ends with the same value, we can replace that value with a constant.


## Constant propagation (example)

$$
\begin{aligned}
& \text { if (i > 3) \{ } \\
& \text { a = 2; } \\
& \text { b = 6; } \\
& \text { \} else \{ } \\
& \mathrm{a}=4 \text {; } \\
& \text { b = 4; } \\
& \text { \} } \\
& c=b+a ;
\end{aligned}
$$



## Constant Propagation

- Need for generalization of flow functions (change of variables not known in advance, e.g. user input)
- Special lattice
- If given variable can have more values $\Rightarrow$ join $\Rightarrow$ go down in lattice

$x=$ nota constant


## Constant propagation example

( $a, b, c, d$ )


## Constant propagation example



## The best solution

- To get the best solution we need to compute flow functions for all paths in the program (MOP = meet over paths)
- In case of loop there are infinitely many paths $\Rightarrow$ not computable
- Instead we compute with edges between BBs (MFP = maximum fixpoint)


## References

围 T. Vojnar, L. Holik Formal Analysis and Verification, lecture 10. 2011/2012.

