## A New Interpretation of the Decipherability

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## Abstract

We define the "quasi code" H as follows: Let  $\Sigma$  and  $\Delta$  be two finite alphabets. Denote H a finite subset of  $2^{\Delta^+} \setminus \emptyset$ . We define the function  $\overline{f}: \Sigma \to H$ , where  $\overline{f}$  is called "quasi coding" of  $\Sigma$ . A quasi code H is called decipherable if, whenever  $f(x_1), \ldots, f(x_n), f(y_1), \ldots, f(y_m)$  are in H and satisfy  $f(x_1) \ldots f(x_n) = f(y_1) \ldots f(y_m)$ , then n = m and  $f(x_i) = f(y_i)$  for all  $i, 1 \leq i \leq n$ .